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Article

Probabilistic Unawareness

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Abstract: The modeling of awareness and unawareness is a significant topic in the doxastic logic literature, where it is usually tackled in terms of full belief operators. The present paper aims at a treatment in terms of partial belief operators. It draws upon the modal probabilistic logic that was introduced by Aumann (1999) at the semantic level, and then axiomatized by Heifetz and Mongin (2001). The paper embodies in this framework those properties of unawareness that have been highlighted in the seminal paper by Modica and Rustichini (1999). Their paper deals with full belief, but we argue that the properties in question also apply to partial belief. Our main result is a (soundness and) completeness theorem that reunites the two strands—modal and probabilistic—of doxastic logic.

Keywords: unawareness; epistemic logic; probabilistic logic

JEL Classification: C70; C72; D80

1. Introduction

Full (or categorical) beliefs are doxastic attitudes, like those ascribed when one says:

Pierre believes that ϕ .

Modal logic provides a way of modeling full beliefs. It is well known that it suffers from two main cognitive idealizations. The first one is logical omniscience: a family of properties such as the closure of beliefs under logical consequence (from the premise that ϕ implies ψ , infer that $B\phi$ implies $B\psi$, also known as the rule of monotonicity) or substitutability of logically equivalent formulas (from the premise that ϕ is equivalent to ψ , infer that $B\phi$ is equivalent to $B\psi$, also known as the rule of equivalence). The second cognitive idealization is full awareness, which is more difficult to characterize precisely. As a first approximation, let us say that, according to this assumption, the agent is supposed to have a full understanding of the underlying space of possibilities and of the propositions that can be built upon them.

Logicians and computer scientists have devoted much attention to the weakening of logical omniscience. This literature is surveyed in [1], and in particular the two main extant solutions: structures with subjective, logically-impossible states¹ and awareness[©] structures introduced by R. Fagin and J. Halpern [4]². The very same awareness[©] structures are used to weaken the full awareness assumption. More recently, game theorists have become interested in weakening awareness in epistemic and doxastic logic and in related formalisms ([5–10])³. For a recent and detailed survey of models of unawareness, see [14].

¹ See [2,3].

² In the whole paper we use “awareness[©]” to denote the model of [4], to be distinguished from the attitude of awareness.

³ There is also a literature that studies games with unaware players. See, for instance, [11–13].

Doxastic logic is a rather coarse-grained model of doxastic attitudes, because it excludes partial beliefs, i.e., the fact that an agent believes that it is unlikely that ϕ or very likely that ψ . The main formalism for partial beliefs makes use of probabilities in their subjective or epistemic interpretation, where probability values stand for degrees of belief. There is a noteworthy contrast between modal doxastic logic and those probabilistic models: whereas the former make beliefs explicit (part of the formal language), they are left implicit in the latter. However, one may enrich the syntax with explicit partial belief operators. For instance, R. Aumann has introduced in [15] an operator $L_\alpha\phi$ interpretable as

the agent believes at least to degree α that ϕ

A possible-world semantics is given for these operators (which is inspired by [16])⁴. This semantics has been axiomatized by [19] under the form of a weak (soundness and) completeness theorem. This probabilistic logic is the true counterpart of Kripkean epistemic logic for degrees of beliefs, and it is the framework of this paper.

This probabilistic logic suffers from the same cognitive idealizations as doxastic logic: logical omniscience and full awareness. In a preceding paper [20], we dealt with the problem of logical omniscience in probabilistic logic. Our proposal was mainly based on the use of so-called impossible states, i.e., subjective states where the logical connectives can have a non-classical behavior. The aim of the present paper is to enrich probabilistic logic with a modal logic of unawareness. Our main proposal is a generalization of Aumann's semantics that uses impossible states like those of [5] and provably satisfies a list of intuitive requirements. Our main result is a weak completeness theorem like the one demonstrated by [19], but adapted to the richer framework that includes awareness. To our knowledge, [21] is the closest work to ours: in this paper, the state-space model with "interactive unawareness" introduced in [8] is extended to probabilistic beliefs in order to deal with the issue of speculative trade. One of the differences with our paper is that their framework is purely set-theoretical, whereas we rely on a formal language.

The remainder of the paper proceeds as follows. In Section 2, we try to provide some intuitions about the target attitudes, awareness and unawareness. Section 3 presents briefly probabilistic logic, and notably the axiom system of [19] (that will be called 'system HM '). In Section 4, we vindicate a slightly modified version of the Generalized Standard Structures of [5]. Section 5 contains the main contribution of the paper: a logic for dealing with unawareness in probabilistic logic⁵. Our axiom system (named 'system HM_U ') enriches the probabilistic logic with an awareness operator and accompanying axioms. Section 6 concludes.

2. Awareness and Unawareness

2.1. Some Intuitions

Unawareness is a more elusive concept than logical omniscience. This section gives insights on the target phenomena and puts forward properties that a satisfactory logic of (un)awareness should embody. Following the lead of [5], we may say that there is unawareness when

- there is "ignorance about the state space"
- "some of the facts that determine which state of nature occurs are not present in the subject's mind"
- "the agent does not know, does not know that she does not know, does not know that she does not know that she does not know, and so on..."

⁴ This is not the only way to proceed. Fagin, Halpern and Moses introduced in [17] an operator $w(\phi)$ and formulas $a_1w(\phi_1) + \dots + a_nw(\phi_n) \geq c$ interpretable as "the sum of a_1 times the degree of belief in ϕ_1 and...and a_n times the degree of belief in ϕ_n is at least as great as c ". For a recent survey of probabilistic logic, see [18].

⁵ B. Schipper in [22] pointed out that a similar result has been stated in an unpublished paper by T. Sadzik; see [23]. The framework is slightly different from ours.

Here is an illustrative example. Pierre plans to rent a house for the next holiday, and from the observer's point of view, there are three main factors relevant to his choice:

- p : the house is no more than 1 km far from the sea
- q : the house is no more than 1 km far from a bar
- r : the house is no more than 1 km far from an airport

There is an intuitive distinction between the two following doxastic states:

State (i): Pierre is undecided about r 's truth: he neither believes that r , nor believes that $\neg r$; there are both r -states and $\neg r$ -states that are epistemically accessible to him.

State (ii): the possibility that r does not come up to Pierre's mind. Pierre does not ask himself: 'is there an airport no more than 1 km far from the house?'.

The contrast between the two epistemic states can be rendered in terms of a state space with either a fine or coarse grain⁶. The observer's set of possible states is: where each state is labeled by the sequence of literals that are true in it. This state is also Pierre's in doxastic State (i). The doxastic State (ii), on the other hand, is: Some states in the initial state space have been fused with each other; those that differ only in the truth value they assign to the formula the agent is unaware of, namely r .⁷

2.2. Some Principles in Epistemic Logic

More theoretically, what properties should one expect awareness to satisfy? In what follows:

- $B\phi$ means "the agent believes that ϕ ,"
- $A\phi$ means "the agent is aware that ϕ ".

Here is a list of plausible properties for the operators B and A :

$A\phi \leftrightarrow A\neg\phi$	(symmetry)
$A(\phi \wedge \psi) \leftrightarrow A\phi \wedge A\psi$	(distributivity over \wedge)
$A\phi \leftrightarrow AA\phi$	(self-reflection)
$\neg A\phi \rightarrow \neg A\neg A\phi$	(U-introspection)
$\neg A\phi \rightarrow \neg B\phi \wedge \neg B\neg B\phi$	(plausibility)
$\neg A\phi \rightarrow (\neg B)^n \phi \forall n \in \mathbb{N}$	(strong plausibility)
$\neg B\neg A\phi$	(BU-introspection)

Natural as they are, these properties cannot be jointly satisfied in Kripkean doxastic logic. This has been recognized by [6], who show that it is impossible to have both:

- a non-trivial awareness operator that satisfies plausibility, U-introspection and BU-introspection and
- a belief operator that satisfies either necessitation or the rule of monotonicity⁸.

⁶ This is close to the "small world" concept of [24]. In Savage's language, "world" means the state space or set of possible worlds, itself.

⁷ Once again, the idea is already present in [24]: "...a smaller world is derived from a larger by neglecting some distinctions between states". The idea of capturing unawareness with the help of coarse-grained or subjective state spaces is widely shared in the literature; see, for instance, [8] or [9]. By contrast, in the framework of first-order epistemic logic, unawareness is construed as unawareness of some objects in the domain of interpretation by [25]. This approach is compared with those based on subjective state spaces in [26].

⁸ For a recent elaboration on the results of [6], see [27].

Of course, the standard belief operator of epistemic logic does satisfy both necessitation and the rule of monotonicity. The main challenge is therefore to build a logic of belief and awareness that supports the above intuitive principles. Since necessitation and the rule of monotonicity are nothing but forms of logical omniscience, it becomes a major prerequisite to weaken the latter. Indeed, both the generalized standard structures of [5] and the awareness[©] structures of [4] do weaken logical omniscience.

2.3. Some Principles in Probabilistic Logic

Probabilistic logics are both lesser known than and not so well unified as modal doxastic logics. The syntactic framework, in particular, varies from one to another. The logic on which this paper is based relies on a language that is quite similar to that of doxastic logic, and it therefore can be seen as constituting a probabilistic modal logic: its primary doxastic operators are L_a , where a is a rational number between zero and one (“the agent believes at least to degree a that...”)⁹ We can express the relevant intuitive principles for L_a as:

$A\phi \leftrightarrow A\neg\phi$	(symmetry)
$A(\phi \wedge \psi) \leftrightarrow A\phi \wedge A\psi$	(distributivity over \wedge)
$A\phi \leftrightarrow AA\phi$	(self-reflection)
$\neg A\phi \rightarrow \neg A\neg A\phi$	(U-introspection)
$\neg A\phi \rightarrow \neg L_a\phi \wedge \neg L_a\neg L_a\phi$	(plausibility)
$\neg A\phi \rightarrow (\neg L_a)^n\phi \forall n \in \mathbb{N}$	(strong plausibility)
$\neg L_a\neg A\phi$	(L_a U-introspection)
$L_0\phi \leftrightarrow A\phi$	(minimality)

Seven of these eight principles are direct counterparts of those put forward for modal doxastic logic, minimality being the exception. On the one hand, if an agent believes to some degree (however small) that ϕ , then he or she is aware of ϕ . This conditional is intuitive for a judgmental rather than for a purely behavioral conception of partial beliefs, according to which degrees of beliefs are causal determinants of behavior, which may or may not be consciously grasped by the agent¹⁰. The reverse conditional roughly means that an agent aware of ϕ has some degree of belief toward ϕ . This directly echoes Bayesian epistemology. These eight principles may be seen as a set of requirements for a satisfactory probabilistic logic¹¹.

3. Probabilistic (Modal) Logic

This section briefly reviews the main concepts of probabilistic logic following [15,19]¹². Probabilistic logic is of course related to the familiar models of beliefs where the doxastic states are represented by a probability distribution on a state space (or on the formulas of a propositional language), but the doxastic operator is made explicit here. Syntactically, as we have already said, this means that the language is endowed with a family of operators L_a . Semantically, there are sets of possible states (or “events”) corresponding to the fact that an agent believes (or does not believe) such and such formula at least to such and such degree.

3.1. Language

Definition 1 (probabilistic language). *The set of formulas of a probabilistic language $\mathcal{L}^L(At)$ based on a set At of propositional variables is defined by:*

$$\phi ::= p \mid \perp \mid \top \mid \neg\phi \mid \phi \wedge \psi \mid L_a\phi$$

⁹ By contrast, in [28,29], one considers formulas like $a_1w(\phi_1) + \dots + a_nw(\phi_n) \geq b$, where a_1, \dots, a_n, b are integers, ϕ_1, \dots, ϕ_n are propositional formulas and $w(\phi)$ is to be interpreted as the probability of ϕ .

¹⁰ For a philosophical elaboration on this distinction, see [30].

¹¹ A similar list of properties is proven in [21].

¹² Economists are leading contributors to the study of explicit probabilistic structures because they correspond to the so-called type spaces, which are basic to games of incomplete information since [16]. See [31].

where $p \in At$ and $a \in [0, 1] \cap \mathbb{Q}$.

From this, one may define two derived belief operators:

- $M_a\phi = L_{1-a}\neg\phi$ (the agent believes at most to degree a that ϕ)
- $E_a\phi = M_a\phi \wedge L_a\phi$ (the agent believes exactly to degree a that ϕ)¹³

3.2. Semantics

Probabilistic structures (PS), as introduced by [15], aim at interpreting the formal language we just defined. They are the true probabilistic counterpart of Kripke structures for epistemic logic. In particular, iterated beliefs are allowed because a probability distribution is attributed to each possible state, very much like a Kripkean accessibility relation. We follow the definition of [19]:

Definition 2 (probabilistic structures). A probabilistic structure for $\mathcal{L}^L(At)$ is a four-tuple $\mathcal{M} = (S, \Sigma, \pi, P)$ where:

- (i) S is a state space
- (ii) Σ is a σ -field of subsets of S
- (iii) $\pi : S \times At \rightarrow \{0, 1\}$ is a valuation for S s.t. $\pi(\cdot, p)$ is measurable for every $p \in At$
- (iv) $P : S \rightarrow \Delta(S, \Sigma)$ is a measurable mapping from S to the set of probability measures on Σ endowed with the σ -field generated by the sets

$$\{\mu \in \Delta(S, \Sigma) : \mu(E) \geq a\} \forall E \in \Sigma, a \in [0, 1].$$

Definition 3. The satisfaction relation, labeled \models , extends π to every formula of the language according to the following conditions:

- (i) $\mathcal{M}, s \models p$ iff $\pi(p, s) = 1$
- (ii) $\mathcal{M}, s \models \phi \wedge \psi$ iff $\mathcal{M}, s \models \phi$ and $\mathcal{M}, s \models \psi$
- (iii) $\mathcal{M}, s \models \neg\phi$ iff $\mathcal{M}, s \not\models \phi$
- (iv) $\mathcal{M}, s \models L_a\phi$ iff $P(s)(\llbracket\phi\rrbracket) \geq a$

As usual, $\llbracket\phi\rrbracket$ denotes the set of states where ϕ is true, or the proposition expressed by ϕ . From a logical point of view, one of the most striking features of probabilistic structures is that compactness does not hold. Let $\Gamma = \{L_{1/2-1/n}\phi : n \geq 2, n \in \mathbb{N}\}$ and $\psi = \neg L_{1/2}\phi$. For each finite $\Gamma' \subset \Gamma$, $\Gamma' \cup \{\psi\}$ is satisfiable, but $\Gamma \cup \{\psi\}$ is not. As a consequence, an axiomatization of probabilistic structures will provide at best a weak completeness theorem.

3.3. Axiomatization

Explicit probabilistic structures were not axiomatized in [15]. To deal with this issue, an axiom system was proposed in [19] that is (weakly) complete for these structures. We coin it the system HM^{14} .

¹³ Since a is a rational number and the structures will typically include real-valued probability distributions, it may happen in some state that for no a , it is true that $E_a\phi$. It happens when the probability assigned to ϕ is a real, but non-rational number.

¹⁴ The work in [19] calls this system $\Sigma+$.

System <i>HM</i>
(PROP) Instances of propositional tautologies
(MP) From ϕ and $\phi \rightarrow \psi$, infer ψ
(L1) $L_0\phi$
(L2) $L_a\top$
(L3) $L_a\phi \rightarrow \neg L_b\neg\phi$ ($a + b > 1$)
(L4) $\neg L_a\phi \rightarrow M_a\phi$
(DefM) $M_a\phi \leftrightarrow L_{1-a}\neg\phi$
(RE) From $\phi \leftrightarrow \psi$ infer $L_a\phi \leftrightarrow L_a\psi$
(B) From $((\phi_1, \dots, \phi_m) \leftrightarrow (\psi_1, \dots, \psi_n))$ infer $((\bigwedge_{i=1}^m L_{a_i}\phi_i) \wedge (\bigwedge_{j=2}^n M_{b_j}\psi_j) \rightarrow L_{(a_1+\dots+a_m)-(b_1+\dots+b_n)}\psi_1)$

The inference rule (B) deserves attention. The content and origin of (B) is explained in [19], so we can be brief. The pseudo-formula $((\phi_1, \dots, \phi_m) \leftrightarrow (\psi_1, \dots, \psi_n))$ is an abbreviation for:

$$\bigwedge_{k=1}^{\max(m,n)} \phi^{(k)} \leftrightarrow \psi^{(k)}$$

where:

$$\phi^{(k)} = \bigvee_{1 \leq l_1 < \dots < l_k \leq m} (\phi_{l_1} \wedge \dots \wedge \phi_{l_k})$$

(if $k > m$, by convention $\phi^{(k)} = \perp$). Intuitively, $\phi^{(k)}$ “says” that at least k of the formulas ϕ_i is true. The meaning of (B) is simpler to grasp when it is interpreted set-theoretically. Associate (E_1, \dots, E_m) and (F_1, \dots, F_n) , two sequences of events, with the sequences of formulas (ϕ_1, \dots, ϕ_m) and (ψ_1, \dots, ψ_n) . Then, the premise of (B) is a syntactical rendering of the idea that the sum of the characteristic functions is equal, i.e., $\sum_{i=1}^m \mathbf{1}_{E_i} = \sum_{j=1}^n \mathbf{1}_{F_j}$. If $P(E_i) \geq \alpha_i$ for $i = 1, \dots, n$ and $P(F_j) \leq \beta_j$ for $j = 2, \dots, m$, then $P(F_1)$ has to “compensate”, i.e.,

$$P(F_1) \geq (\alpha_1 + \dots + \alpha_n) - (\beta_2 + \dots + \beta_m)$$

The conclusion of (B) is a translation of this “compensation”. It is very powerful from the probabilistic point of view and plays a crucial role in the (sophisticated) completeness proof. In comparison with the modal doxastic logic, one of the issues is that it is not easy to adapt the usual proof method, i.e., that of canonical models. More precisely, with Kripke logics, there is a natural accessibility relation on the canonical state space. Here, we need to prove the existence of a canonical probability distribution from relevant mathematical principles. This step is linked to a difficulty that is well known in the axiomatic study of quantitative and qualitative probability: how to ensure a (numerical) probabilistic representation for a finite structure of qualitative probability. A principle similar to (B) has been introduced in an set-theoretical axiomatic framework by [32] and imported into a probabilistic modal logic (with a qualitative binary operator) by [33].

4. A Detour by Doxastic Logic without Full Awareness

Our probabilistic logic without full awareness is largely an adaptation of the generalized standard structures (GSS) introduced by [5] to deal with unawareness in doxastic logic. Actually, we will slightly modify the semantics of [5] and obtain a partial semantics for unawareness. Before giving our own logical system, we remind the reader of GSSs of the case of doxastic logic.

4.1. Basic Heuristics

Going back to the motivating example, suppose that the “objective” state space is based on the set of atomic formulas $At = \{p, q, r\}$ as in Figure 1. Suppose furthermore that the actual state is $s = pqr$ and that in this state, Pierre believes that p , is undecided about q and is unaware of r . In the language

used in [5], one would say that the actual state is “projected” to a subjective state $\rho(s) = pq$ of a “subjective” state space based on the set of atomic formulas that the agent is aware of, i.e., p and q (Figure 2). In Kripkean doxastic logic, the agent’s accessibility relation selects, for each possible state s , the set of states $R(s)$ that are epistemically possible for him or her. GSSs define an accessibility relation on this subjective state space. In Figure 3, projection is represented by a dotted arrow and accessibility by solid arrows.

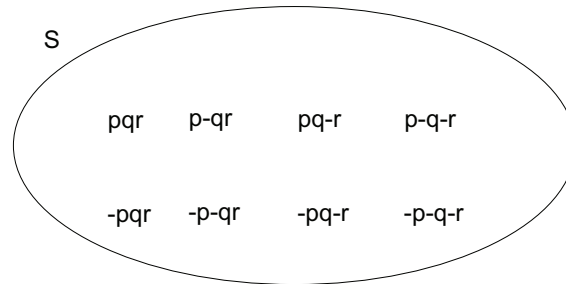


Figure 1. An objective state space.

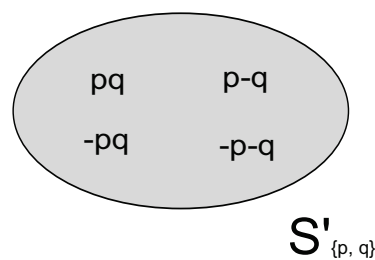


Figure 2. A subjective state space.

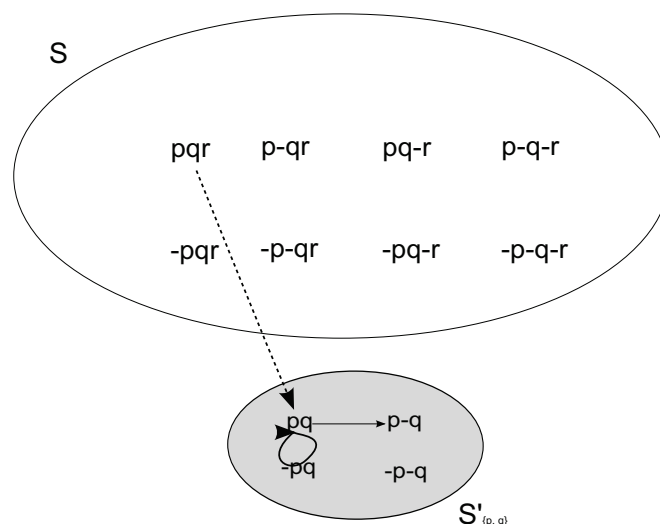


Figure 3. Projection of an objective state in a subjective state space.

This picture does not represent all projections between objective states and subjective states, and it corresponds to only one subjective state space. Generally, there are as many subjective state

spaces S'_X as there are subsets X of the set At on which the objective state space is based¹⁵. It is crucial to specify in the right way the conditions on the projection $\rho(\cdot)$ from objective to subjective states. Suppose that another objective state s' is projected to pq , as well; then, two conditions should be obtained. First, $s = pqr$ and s' should agree on the atomic formulas the agent is aware of; so, for instance, s' could not be $\neg pqr$, since the agent is aware of p . Second, the states accessible from s and s' should be the same. Another natural assumption is that all of the states accessible from a given state are located in the same “subjective” state space (see (v)(2) in Definition 4 below). Figure 4 pictures a GSS more faithfully than the preceding one.

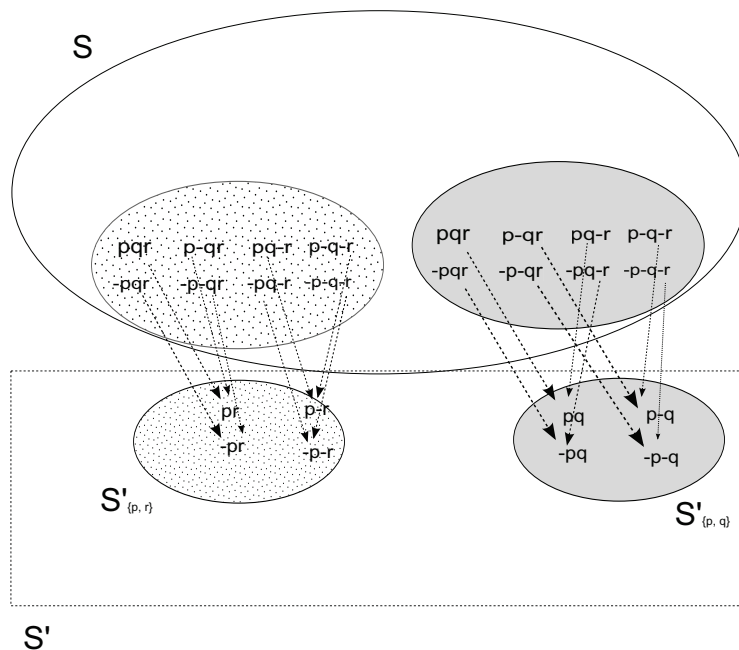


Figure 4. Partial picture of a generalized standard structure (GSS).

4.2. Generalized Standard Structures

The next definition is a precise rendering of the intuitive ideas. We have followed [7] rather than [5], notably because he makes clear that GSSs can be seen as structures with impossible states.

Definition 4. A GSS is a t -tuple $\mathcal{M} = (S, S', \pi, R, \rho)$:

- (i) S is a state space
- (ii) $S' = \bigcup_{X \subseteq At} S'_X$ (where S'_X are disjoint) is a (non-standard) state space
- (iii) $\pi : S \times At \rightarrow \{0, 1\}$ is a valuation for S
- (iv) $R : S \rightarrow \wp(S')$ is an accessibility relation for S
- (v) $\rho : S \rightarrow S'$ is an onto map s.t.

(1) if $\rho(s) = \rho(t) \in S'_X$, then (a) for each atomic formula $p \in X$, $\pi(s, p) = \pi(t, p)$ and (b) $R(s) = R(t)$ and

(2) if $\rho(s) \in S'_X$, then $R(s) \subseteq S'_X$

One can extend R and π to the whole state space:

¹⁵ In an objective state space, the set of states need not reflect the set of possible truth-value assignments to propositional variables (as is the case in our example).

- (vi) $\pi' : S' \times At \rightarrow \{0, 1\}$ is a valuation for S' s.t. for all $X \subseteq At$ for all $s' \in S'_X$, $\pi'(s', p) = 1$ iff (a) $p \in X$ and (b) for all $s \in \rho^{-1}(s')$, $\pi(s, p) = 1$. We note $\pi^* = \pi \cup \pi'$.
- (vii) $R' : S' \rightarrow \wp(S')$ is an accessibility relation for S s.t. for all $X \subseteq At$ for all $s' \in S'_X$, $R'(s') = R(s)$ for some $s \in \rho^{-1}(s')$. We note $R^* = R \cup R'$.

In comparison with Kripke structures, a modification is introduced as regards negation. In a subjective state $s' \in S'_X$, for a negated formula $\neg\phi$ to be true, it has to be not only that ϕ is not true, but also that ϕ belongs to the sub-language induced by X . Semantic partiality follows: it may be the case that in some s' , neither ϕ , nor $\neg\phi$ is true (this is why subjective states are impossible states). The main reason why, following [7], we introduce this semantics is that it is a very simple way of inducing the “right” kind of partiality. This will be shown in the next subsection. In the sequel, $\mathcal{L}^{BA}(X)$ denotes the language containing the operators B (full beliefs) and A (awareness) and based on the set X of propositional variables.

Definition 5. The satisfaction relation for GSS is defined for each $s^* \in S^* = S \cup S'$:

- (i) $\mathcal{M}, s^* \models p$ iff $\pi^*(s^*, p) = 1$
- (ii) $\mathcal{M}, s^* \models \phi \wedge \psi$ iff $\mathcal{M}, s^* \models \phi$ and $\mathcal{M}, s^* \models \psi$
- (iii) $\mathcal{M}, s^* \models \neg\phi$ iff $\mathcal{M}, s^* \not\models \phi$ and either $s^* \in S$, or $s^* \in S'_X$ and $\phi \in \mathcal{L}^{BA}(X)$
- (iv) $\mathcal{M}, s^* \models B\phi$ iff for each $t^* \in R^*(s^*)$, $\mathcal{M}, t^* \models \phi$
- (v) $\mathcal{M}, s^* \models A\phi \Leftrightarrow \mathcal{M}, s^* \models B\phi \vee B\neg B\phi$

Example 1. In Figure 3, let us consider $s = pqr$. $\mathcal{M}, \rho(s) \not\models r$ given Clause (vi) of the definition of GSS. However, $\mathcal{M}, \rho(s) \not\models \neg r$ given Clause (iii) of the definition of the satisfaction relation. Clause (iv) of the satisfaction relation and Clause (v)(2) of the GSS imply that $\mathcal{M}, \rho(s) \not\models Br$ and that $\mathcal{M}, s \not\models Br$. However, Clause (iii) of the satisfaction relation again implies that $\mathcal{M}, \rho(s) \not\models \neg Br$. The same holds in the state accessible from $\rho(s)$. Therefore, $\mathcal{M}, s \not\models B\neg Br$. By Clause (v) of the satisfaction relation, this implies that $\mathcal{M}, s \models \neg Ar$.

4.3. Partial Generalized Standard Structures

The preceding definition characterizes awareness in terms of beliefs: Pierre is unaware of ϕ if, and only if, he does not believe that ϕ and does not believe that he does not believe that ϕ . This is unproblematic when one studies, as [5] do, partitional structures, i.e., structures where the accessibility relation is an equivalence relation and, thus, induces a partition of the state space. Game theorists rely extensively on this special case, which has convenient properties. In this particular case, the fact that an agent does not believe that ϕ and does not believe that she/he does not believe that ϕ implies that she/he does not believe that she/he does not believe...that she/he does not believe that ϕ , at all levels of iteration¹⁶. However, without such an implication, the equivalence between the fact that an agent is unaware of ϕ and the fact that she/he does not believe that ϕ and does not believe that she/he does not believe that ϕ is dubious, at least in one of the two directions.

We therefore need a more general characterization of awareness and unawareness. Our proposal proceeds from the following observation: in the definition of satisfaction for GSS, the truth-conditions for negated formulas introduce (semantic) partiality. If $p \notin X$ and $s^* \in S'_X$, then neither $\mathcal{M}, s^* \models p$ nor $\mathcal{M}, s^* \models \neg p$ obtain. Let us indicate by $\mathcal{M}, s^* \uparrow \phi$ that the formula ϕ is undefined at s^* and by $\mathcal{M}, s^* \downarrow \phi$ that it is defined. The following is true:

Fact 1. Let \mathcal{M} be a GSS and $s^* \in S'_X$ for some $X \subseteq At$. Then:

¹⁶ Intuitively: if the agent does not believe that ϕ , this means that not every state of the relevant partition's cell makes ϕ true. If ϕ belonged to the sub-language associated with the relevant subjective state space, since the accessibility relation is partitional, this would imply that $\neg B\phi$ would be true in every state of the cell. However, by hypothesis, the agent does not believe that she/he does not believe that ϕ . We have therefore to conclude that ϕ does not belong to the sub-language of the subjective state space (see Fact 1 below). Hence, no formula $B\neg B\neg B\dots\neg B\phi$ can be true.

$$\mathcal{M}, s^* \Downarrow \phi \text{ iff } \phi \in \mathcal{L}^{BA}(X)$$

Proof. see Appendix A.2. \square

We suggest to keep the underlying GSS, but change the definition of (un)awareness. Semantical partiality, stressed above, is an attractive guide. In our introductory example, one would like to say that the possible states that Pierre conceives of do not “answer” the question “Is it true that r ?”, whereas they do answer the questions “Is it true that p ?” and “Is it true that q ?”. In other words, the possible states that Pierre conceives of make neither r , nor $\neg r$ true. Awareness can be defined semantically in terms of partiality:

$$\mathcal{M}, s \models A\phi \text{ iff } \mathcal{M}, \rho(s) \Downarrow \phi$$

Of course, the appeal of this characterization depends on the already given condition: if $\rho(s) \in S'_X$, then $R(s) \subseteq S'_X$. Let us call a partial GSS a GSS where the truth conditions of the (un)awareness operator are in terms of partiality:

Fact 2. *Symmetry, distributivity over \wedge , self-reflection, U-introspection, plausibility and strong plausibility are valid under partial GSSs. Furthermore, BU-introspection is valid under serial partial GSS.*

Proof. This is left to the reader. \square

5. Probabilistic Logic without Full Awareness

5.1. Language

Definition 6 (Probabilistic language with awareness). *The set of formulas of a probabilistic language with awareness $\mathcal{L}^{LA}(At)$ based on a set At of propositional variables is defined by:*

$$\phi ::= p \mid \perp \mid \top \mid \neg\phi \mid \phi \wedge \psi \mid L_a\phi \mid A\phi$$

where $p \in At$ and $a \in [0, 1] \cap \mathbb{Q}$.

5.2. Generalized Standard Probabilistic Structures

Probabilistic structures make a full awareness assumption, exactly in the same way that Kripke structures do. An obvious way to weaken this assumption is to introduce in the probabilistic setting the same kind of modification as the one investigated in the previous section. The probabilistic counterpart of generalized standard Structures are the following generalized standard probabilistic structures (GSPS):

Definition 7 (Generalized standard probabilistic structure). *A generalized standard probabilistic structure for $\mathcal{L}^{LA}(At)$ is a t -tuple*

$$\mathcal{M} = (S, (S'_X)_{X \subseteq At}, (\Sigma'_X)_{X \subseteq At}, \pi, (P'_X)_{X \subseteq At}) \text{ where:}$$

- (i) S is a state space.
- (ii) S'_X where $X \subseteq At$ are disjoint “subjective” state spaces. Let $S' = \bigcup_{X \subseteq At} S'_X$.
- (iii) For each $X \subseteq At$, Σ'_X is a σ -field of subsets of S'_X .
- (iv) $\pi : S \times At \rightarrow \{0, 1\}$ is a valuation.
- (v) $P'_X : S'_X \rightarrow \Delta(S'_X, \Sigma'_X)$ is a measurable mapping from S'_X to the set of probability measures on Σ'_X endowed with the σ -field generated by the sets $\{\mu \in \Delta(S'_X, \Sigma'_X) : \mu(E) \geq a\}$ for all $E \in \Sigma'_X, a \in [0, 1]$.
- (vi) $\rho : S \rightarrow S'$ is an onto map s.t. if $\rho(s) = \rho(t) \in S'_X$, then for each atomic formula $p \in X$, $\pi(s, p) = \pi(t, p)$.
By definition, $P^*(s^*) = P^*(\rho(s^*))$ if $s \in S$ and $P^*(s^*) = P'_X(s^*)$ if $s^* \in S'_X$.

(vi) $\pi' : S' \times At \rightarrow \{0, 1\}$ extends π to S' as follows: for all $s' \in S'_X$, $\pi'(s', p) = 1$ iff $p \in X$ and for all $s \in \rho^{-1}(s')$, $\pi(s, p) = 1$ ¹⁷. For every $p \in At$, $s' \in S'_X$, $\pi'(\cdot, p)$ is measurable w.r.t. (S'_X, Σ'_X) .

Two comments are in order. First, Clause (iv) does not introduce any special measurability condition on the newly-introduced awareness operator (by contrast, there is still a condition for the doxastic operators). The reason is that in a given subjective state space S'_X , for any formula ϕ , either there is awareness of ϕ everywhere, and in this case $[[A\phi]] \cap S'_X = S'_X$, or there is never awareness of ϕ , and in this case $[[A\phi]] \cap S'_X = \emptyset$. These two events are of course already in any Σ'_X . Second, Clause (v) imposes conditions on the projection ρ . With respect to GSSs, the only change is that we do not require something like: if $\rho(s) \in S'_X$, then $R(s) \subseteq S'_X$. The counterpart would be that if $\rho(s) \in S'_X$, then $Supp(P(s)) \subseteq S'_X$. However, this is automatically satisfied by the definition.

Example 2. In Figure 5, the support of the probability distribution associated with $s = pqr$ is $\{\rho(s) = pq, p\bar{q}\}$, $P'_{\{p,q\}}(\rho(s))(pq) = a$ and $P'_{\{p,q\}}(\rho(s))(p\bar{q}) = 1 - a$.

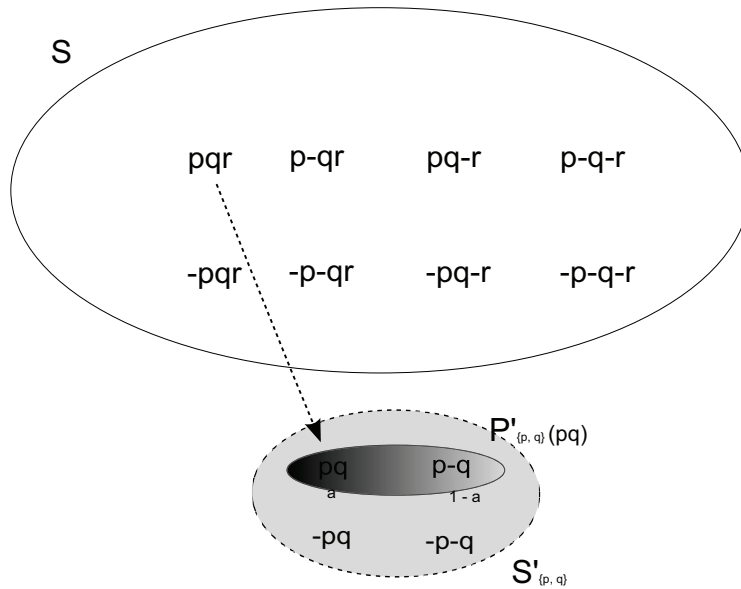


Figure 5. Projection of an objective state to a subjective state space in a probabilistic setting.

Definition 8 (Satisfaction relation for GSPS). *The satisfaction relation for GSPS is defined for each $s^* \in S^* = S \cup S'$:*

- (i) $\mathcal{M}, s^* \models p$ iff $\pi(s^*, p) = 1$
- (ii) $\mathcal{M}, s^* \models \phi \wedge \psi$ iff $\mathcal{M}, s^* \models \phi$ and $\mathcal{M}, s^* \models \psi$
- (iii) $\mathcal{M}, s^* \models \neg\phi$ iff $\mathcal{M}, s^* \not\models \phi$ and either $s^* \in S$, or $s^* \in S'_X$ and $\phi \in \mathcal{L}^{LA}(X)$
- (iv) $\mathcal{M}, s^* \models L_a\phi$ iff $P^*(s^*)([[\phi]]) \geq a$ and $\mathcal{M}, \rho(s) \Downarrow \phi$
- (v) $\mathcal{M}, s^* \models A\phi$ iff $\mathcal{M}, \rho(s^*) \Downarrow \phi$

The following fact is the counterpart for GSPS of what was proven above for GSS.

Fact 3. *Let \mathcal{M} be a GSPS and $s' = \rho(s) \in S'_X$ for some $s \in S$. Then:*

$$\mathcal{M}, s' \Downarrow \phi \text{ iff } \phi \in \mathcal{L}^{LA}(X)$$

¹⁷ It follows from Clause (v) above that this extension is well defined.

Proof. The proof is analogous to the one provided for Fact 1 above. \square

One can show that all of the properties mentioned in Section 2.3 are valid under GSPSs.

Proposition 1. For all GSPS \mathcal{M} and all standard states $s \in S$, the following formulas are satisfied:

$A\phi \leftrightarrow A\neg\phi$	(symmetry)
$A(\phi \wedge \psi) \leftrightarrow A\phi \wedge A\psi$	(distributivity over \wedge)
$A\phi \leftrightarrow AA\phi$	(self-reflection)
$\neg A\phi \rightarrow \neg A\neg A\phi$	(U-introspection)
$\neg A\phi \rightarrow \neg L_a\phi \wedge \neg L_a\neg L_a\phi$	(plausibility)
$\neg A\phi \rightarrow (\neg L_a)^n\phi \forall n \in \mathbb{N}$	(strong plausibility)
$\neg L_a\neg A\phi$	(L_a U-introspection)
$L_0\phi \leftrightarrow A\phi$	(minimality)

Proof. This is left to the reader. \square

5.3. Axiomatization

Proposition 1 suggests that GSPSs provide a plausible analysis of awareness and unawareness in a probabilistic setting. To have a more comprehensive understanding of this model, we need to investigate its logical properties. It turns out that an axiom system can be given that is weakly complete with respect to GSPS. We call it system HM_U ¹⁸.

System HM_U
(PROP) Instances of propositional tautologies
(MP) From ϕ and $\phi \rightarrow \psi$, infer ψ
(A1) $A\phi \leftrightarrow A\neg\phi$
(A2) $A(\phi \wedge \psi) \leftrightarrow A\phi \wedge A\psi$
(A3) $A\phi \leftrightarrow AA\phi$
(A4 _L) $A\phi \leftrightarrow AL_a\phi$
(A5 _L) $A\phi \rightarrow L_1A\phi$
(L1 _U) $A\phi \leftrightarrow L_0\phi$
(L2 _U) $A\phi \rightarrow L_a(\phi \vee \neg\phi)$
(L3) $L_a\phi \rightarrow \neg L_b\neg\phi$ ($a + b > 1$)
(L4 _U) $(\neg L_a\phi \wedge A\phi) \rightarrow M_a\phi$
(RE _U) From $\phi \leftrightarrow \psi$ and $Var(\phi) = Var(\psi)$, infer $(L_a\phi \leftrightarrow L_a\psi)$
(B _U) From $((\phi_1, \dots, \phi_m) \leftrightarrow (\psi_1, \dots, \psi_n))$, infer: $((\bigwedge_{i=1}^m L_{a_i}\phi_i) \wedge (\bigwedge_{j=2}^n M_{b_j}\psi_j)) \rightarrow (A\psi_1 \rightarrow L_{(a_1+\dots+a_m)-(b_1+\dots+b_n)}\psi_1)$

Some comments are in order. (a) Axioms (A1)–(A5_L) concern the awareness operator and its relationship with the doxastic operators. The subscript “L” indicates an axiom that involves a probabilistic doxastic operator, to be distinguished from its epistemic counterpart indicated by “B” in the axiom system for doxastic logic reproduced in Appendix A.3. The other axioms and inference rules were roughly part of probabilistic logic without the awareness operator appearing in them. Subscript “U” indicates a modification with respect to the System HM due to the presence of the

¹⁸ In what follows, $Var(\phi)$ denotes the set of propositional variables occurring in ϕ .

awareness operator. (b) Axiom ($L1_U$) substitutes $A\phi \leftrightarrow L_0\phi$ for $L_0\phi$. This means that in our system, the awareness operator can be defined in terms of the probabilistic operator. However, this does not imply that a non-standard semantics is not needed: if one defines $A\phi$ as $L_0\phi$, $A\phi$ is valid for any ϕ in standard semantics¹⁹. (c) The relationship between logical omniscience and full awareness comes out clearer in the restriction we have to impose on the rule of equivalence (RE)²⁰. It is easy to see why in semantical terms why the rule of equivalence no longer holds universally. Suppose that $L_{1/2}p$ holds at some state of some model. From propositional logic, we know that $p \equiv (p \wedge q) \vee (p \wedge \neg q)$. However, if the agent is not aware of q , it is not true that $L_{1/2}((p \wedge q) \vee (p \wedge \neg q))$. (d) For the same kind of reason, the inference rule (B) no longer holds universally. Consider for instance $\phi_1 = (p \vee \neg p)$, $\phi_2 = (r \vee \neg r)$, $\psi_1 = (q \vee \neg q)$ and $\psi_2 = (p \vee \neg p)$. Additionally, suppose that the agent is aware of p and r , but not of q . Clearly, the premise of (B), i.e., $((\phi_1, \phi_2) \leftrightarrow (\psi_1, \psi_2))$ is satisfied. Furthermore, the antecedent of (B)'s conclusion, i.e., $L_1\phi_1$ and $M_1\psi_2$ is satisfied, as well. However, since the agent is unaware of q , we cannot conclude what we should conclude were (B) valid, i.e., that $L_1\psi_1$.

We are now ready to formulate our main result.

Theorem 1 (Soundness and completeness of HM_U). *Let $\phi \in \mathcal{L}^{LA}(At)$. Then:*

$$\models_{GSPS} \phi \text{ iff } \vdash_{HM_U} \phi$$

Proof. See the Appendix B. \square

6. Conclusions

This study of unawareness in probabilistic logic could be unfolded later in several directions. First, we did not deal with the extension to a multi-agent framework, an issue tackled recently by [8]; second, we did not investigate applications to decision theory or game theory. However, we would like to end by stressing another issue that is less often evoked and nonetheless conceptually very challenging: the dynamics of awareness. The current framework, as much of the existing work, is static, i.e., it captures the awareness and doxastic states at a given time. It does not tell anything of the fact that, during an inquiry, an agent may become aware of some new possibilities²¹.

Let us consider our initial example where Pierre is aware of p and q , but not of r , and let us suppose that Pierre's partial beliefs are represented by some probability distribution on a subjective state space $S_{\{p,q\}}$. Assume that at some time, Pierre becomes aware of r ; for instance, someone has asked him whether he thinks that r is likely or not. It seems that our framework can be extended to accommodate the situation: Pierre's new doxastic state will be represented on a state space $S_{\{p,q,r\}}$ appropriately connected to the initial one $S_{\{p,q\}}$ (see Figure 6). Typically, a state $s' = pq$ will be split into two fine-grained states $s_1 = pqr$ and $s_2 = pq\neg r$. However, how should Pierre's partial beliefs evolve? Obviously, a naive Laplacian rule according to which the probability assigned to s' is equally allocated to s_1 , and s_2 will not be satisfactory. Are there rationality constraints capable of determining a new probability distribution on $S_{\{p,q,r\}}$? Or should we represent the new doxastic state of the agent by a set of probability distributions?²² We leave the answers to these questions for future investigation.

¹⁹ I thank an anonymous referee for suggesting to me to clarify this point.

²⁰ This restricted rule is reminiscent of the rule RE_{sq} in [5].

²¹ Note that he/she could become unaware of some possibilities as well, but we will not say anything about that.

²² The dynamics of awareness has been studied by [34] in doxastic logic and by [35] in doxastic logic. See also [36].

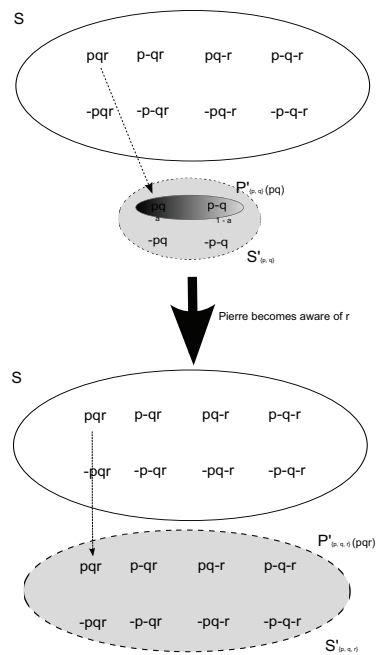


Figure 6. The issue of becoming aware.

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Appendix A. Doxastic Logic: Proofs and Illustrations

Appendix A.1. Illustration

Let us consider a GSS \mathcal{M} where:

- the actual state is $s \in S$
- s is projected in $s_1 \in S'_X$ for some $X \subseteq At$
- $R(s_1) = \{s_2, s_3\}$, $R(s_2) = \{s_2\}$ and $R(s_3) = \{s_3\}$
- $\mathcal{M}, s_2 \models p$ and $\mathcal{M}, s_3 \models \neg p$

The relevant part of the model is represented in Figure A1.

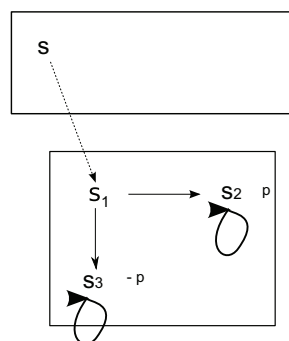


Figure A1. Unawareness without partitions.

It is easy to check that $\mathcal{M}, s \models Up$ and $\mathcal{M}, s \models B(B\neg p \vee Bp)$, since $\mathcal{M}, s_2 \models Bp$ and $\mathcal{M}, s_3 \models B\neg p$.

Appendix A.2. Proof of Fact 1

For greater convenience, we give the proof for partial GSSs, but Fact 1 holds for original GSSs, as well. We have to show that if \mathcal{M} is a partial GSS and $s^* \in S'_X$ for some $X \subseteq At$, then $\mathcal{M}, s^* \Downarrow \phi$ iff $\phi \in \mathcal{L}^{BA}(X)$. The proof is by induction on the complexity of formulas:

- if $\phi := p$, then $\mathcal{M}, s^* \models p$ or $\mathcal{M}, s^* \models \neg p$ iff $(p \in \mathcal{L}^{BA}(X) \text{ and } \pi(\rho^{-1}(s^*), p) = 1)$ or $(p \in \mathcal{L}^{BA}(X) \text{ and not } \pi(\rho^{-1}(s^*), p) = 1)$ iff $p \in \mathcal{L}^{BA}(X)$.
- if $\phi := \neg\psi$, then $\mathcal{M}, s^* \models \phi$ or $\mathcal{M}, s^* \models \neg\phi$ iff $\mathcal{M}, s^* \models \neg\psi$ or $\mathcal{M}, s^* \models \neg\neg\psi$ iff $(\psi \in \mathcal{L}^{BA}(X) \text{ and } \mathcal{M}, s^* \not\models \psi)$ or $(\neg\psi \in \mathcal{L}^{BA}(X) \text{ and } \mathcal{M}, s^* \not\models \neg\psi)$ iff $(\psi \in \mathcal{L}^{BA}(X) \text{ and } \mathcal{M}, s^* \not\models \psi)$ or $(\psi \in \mathcal{L}^{BA}(X) \text{ and } \mathcal{M}, s^* \models \psi)$ iff $\psi \in \mathcal{L}^{BA}(X)$ iff $\neg\psi \in \mathcal{L}^{BA}(X)$
- if $\phi := \psi_1 \wedge \psi_2$, then $\mathcal{M}, s^* \models \phi$ or $\mathcal{M}, s^* \models \neg\phi$ iff $(\mathcal{M}, s^* \models \psi_1 \text{ and } \mathcal{M}, s^* \models \psi_2)$ or $(\psi_1 \wedge \psi_2 \in \mathcal{L}^{BA}(X) \text{ and } (\mathcal{M}, s^* \not\models \psi_1 \text{ or } \mathcal{M}, s^* \not\models \psi_2))$ iff by IH $(\psi_1 \wedge \psi_2 \in \mathcal{L}^{BA}(X) \text{ and } \mathcal{M}, s^* \models \psi_1 \text{ and } \mathcal{M}, s^* \models \psi_2)$ or $(\psi_1 \wedge \psi_2 \in \mathcal{L}^{BA}(X) \text{ and not } (\mathcal{M}, s^* \models \psi_1 \text{ and } \mathcal{M}, s^* \models \psi_2))$ iff $\psi_1 \wedge \psi_2 \in \mathcal{L}^{BA}(X)$
- if $\phi := B\psi$, then $\mathcal{M}, s^* \models \phi$ or $\mathcal{M}, s^* \models \neg\phi$ iff (for each $t^* \in R^*(s^*)$, $\mathcal{M}, t^* \models \phi$) or $(B\psi \in \mathcal{L}^{BA}(X) \text{ and } \mathcal{M}, s^* \not\models B\psi)$ iff, by the induction hypothesis and since each $t^* \in R^*(s^*)$ belongs to $S_X - (B\psi \in \mathcal{L}^{BA}(X) \text{ } \mathcal{M}, s^* \models B\psi)$ or $(B\psi \in \mathcal{L}^{BA}(X) \text{ and } \mathcal{M}, s^* \not\models B\psi)$ iff $B\psi \in \mathcal{L}^{BA}(X)$
- if $\phi := A\psi$, then $\mathcal{M}, s^* \models \phi$ or $\mathcal{M}, s^* \models \neg\phi$ iff $\mathcal{M}, s^* \Downarrow \psi$ or $(A\psi \in \mathcal{L}^{BA}(X) \text{ and } \mathcal{M}, s^* \not\models A\psi)$ iff (by Induction Hypothesis) $\psi \in \mathcal{L}^{BA}(X)$ or $(A\psi \in \mathcal{L}^{BA}(X) \text{ and } \mathcal{M}, s^* \uparrow \psi)$ iff (by the induction hypothesis) $\psi \in \mathcal{L}^{BA}(X)$ or $(A\psi \in \mathcal{L}^{BA}(X) \text{ and } \psi \notin \mathcal{L}^{BA}(X))$ iff $\psi \in \mathcal{L}^{BA}(X)$ iff $A\psi \in \mathcal{L}^{BA}(X)$.

Appendix A.3. An Axiom System for Partial GSSs

We may obtain a complete axiom system for serial partial GSS thanks to [37] who relates GSS and awareness[©] structures. Actually, one obtains a still closer connection with serial partial GSS. Let us first restate the definition of awareness[©] structures.

Definition A1. An awareness[©] structure is a t -tuple (S, π, R, A) where

- (i) S is a state space,
- (ii) $\pi : At \times S \rightarrow \{0, 1\}$ is a valuation,
- (iii) $R \subseteq S \times S$ is an accessibility relation,
- (iv) $A : S \rightarrow \text{Form}(\mathcal{L}^{BA}(At))$ is a function that maps every state in a set of formulas ("awareness[©] set").

The new condition on the satisfaction relation is the following:

$$\mathcal{M}, s \models B\phi \text{ iff } \forall t \in R(s) \mathcal{M}, t \models \phi \text{ and } \phi \in A(s)$$

Let us say that an awareness[©] structure $\mathcal{M} = (S, R, A, \pi)$ is propositionally determined (p.d.) if (1) for each state s , $A(s)$ is generated by some atomic formulas $X \subseteq At$, i.e., $A(s) = \mathcal{L}^{BA}(X)$, and (2) if $t \in R(s)$, then $A(s) = A(t)$.

Proposition A1 (Adapted from Halpern 2001 Theorem 4.1).

1. For every serial p.d. awareness[©] structure \mathcal{M} , there exists a serial partial GSS \mathcal{M}' based on the same state space S and the same valuation π s.t. for all formulas $\phi \in \mathcal{L}^{BA}(At)$ and each possible state s

$$\mathcal{M}, s \models_{a^\circ} \phi \text{ iff } \mathcal{M}', s \models_{p\text{GSS}} \phi$$
2. For every serial partial GSS \mathcal{M} , there exists a serial p.d. awareness[©] structure \mathcal{M}' based on the same state space S and the same valuation π s.t. for all formulas $\phi \in \mathcal{L}^{BA}(At)$ and each possible state s

$$\mathcal{M}, s \models_{p\text{GSS}} \phi \text{ iff } \mathcal{M}', s \models_{a^\circ} \phi$$

An axiom system has been devised in [37] that is (sound and) complete with respect to p.d. awareness[©] structures. An axiom system for serial p.d. awareness[©] structures can be devised by enriching this axiom system with:

$$(D_U) B\phi \rightarrow (\neg B\neg\phi \wedge A\phi)$$

The resulting axiom system, coined KD_U , is this one:

System KD_U
(PROP) Instances of propositional tautologies
(MP) From ϕ and $\phi \rightarrow \psi$, infer ψ
(K) $B\phi \wedge B(\phi \rightarrow \psi) \rightarrow B\psi$
(Gen) From ϕ , infer $A\phi \rightarrow B\phi$
(D _U) $B\phi \rightarrow (\neg B\neg\phi \wedge A\phi)$
(A1) $A\phi \leftrightarrow A\neg\phi$
(A2) $A(\phi \wedge \psi) \leftrightarrow (A\phi \wedge A\psi)$
(A3) $A\phi \leftrightarrow AA\phi$
(A4 _B) $AB\phi \leftrightarrow A\phi$
(A5 _B) $A\phi \rightarrow BA\phi$
(Irr) If no atomic formulas in ϕ appear in ψ , from $U\phi \rightarrow \psi$, infer ψ

The following derives straightforwardly from Proposition 2.

Proposition A2 (Soundness and completeness theorem). *Let $\phi \in \mathcal{L}^{BA}(At)$. Then:*

$$\models_{spGSS} \phi \text{ iff } \vdash_{KD_U} \phi$$

Appendix B. Probabilistic Logic: Proof of the Completeness Theorem for HM_U

Proof. (\Leftarrow). Soundness is easily checked and is left to the reader. (\Rightarrow). We have to show that if $\models_{GSPS} \phi$, then $\vdash_{HM_U} \phi$. The proof relies on the well-known method of filtration. First, we define a restricted language $\mathcal{L}_{[\phi]}^{LA}$ as in [19]: $\mathcal{L}_{[\phi]}^{LA}$ contains:

- as atomic formulas, only $Var(\phi)$, i.e., the atomic formulas occurring in ϕ ,
- only probabilistic operators L_a belonging to the finite set $Q(\phi)$ of rational numbers of the form p/q , where q is the smallest common denominator of indexes occurring in ϕ and
- only formulas of epistemic depth smaller than or equal to that of ϕ (an important point is that we stipulate that the awareness operator A does not add any epistemic depth to a formula: $dp(A\psi) = dp(\psi)$).

As we will show, the resulting language $\mathcal{L}_{[\phi]}^{LA}$ is finitely generated: there is a finite subset \mathfrak{B} of $\mathcal{L}_{[\phi]}^{LA}$ called a base, such that $\forall \psi \in \mathcal{L}_{[\phi]}^{LA}$, there is a formula ψ' in the base, such that $\vdash_{HM_U} \psi \leftrightarrow \psi'$. In probabilistic structures, it is easy to construct such a base²³. The basic idea is this:

- (1) consider D_0 , the set of all the disjunctive normal forms built from $B_0 = Var(\phi)$, the set of propositional variables occurring in ϕ .
- (2) B_k is the set of formulas $L_a\psi$ for all $a \in Q(\phi)$ where ψ is a disjunctive normal form built with “atoms” coming from B_0 to B_{k-1} .
- (3) the construction has to be iterated up to the epistemic depth n of ϕ , hence to B_n . The base \mathfrak{B} is D_n , i.e., the set of disjunctive normal forms built with “atoms” from B_0 to B_n .

²³ The work in [19] leaves the construction implicit.

Obviously, \mathfrak{B} is finite. It can be shown by induction that $\mathcal{L}_{[\phi]}^{LA}$ is finitely generated by \mathfrak{B} . For formulas with a Boolean connective as top connective, this is obvious. For formulas of the form $L_a\psi$, it comes from substitutability under logically-equivalent formulas: by the induction hypothesis, there is a formula ψ' equivalent to ψ in \mathfrak{B} . Therefore, there is in \mathfrak{B} a formula equivalent to $L_a\psi'$. However, since $\vdash_{HM} \psi \leftrightarrow \psi'$, it follows that $\vdash_{HM} L_a\psi \leftrightarrow L_a\psi'$. We will now show how to unfold these ideas formally.

Definition B1. Let $X = \{\psi_1, \dots, \psi_m\}$ be a finite set of formulas. $DNF(X)$ is the set of disjunctive normal forms that can be built from X , i.e., the set of all possible disjunctions of conjunctions of the form $e_1\psi_1 \wedge \dots \wedge e_m\psi_m$ where e_i is a blank or \neg . The members of X are called the atomic components of $DNF(X)$.

Definition B2. The base \mathfrak{B} for a language $\mathcal{L}_{[\phi]}^L$ where $dp(\phi) = n$ is defined as D_n in the following doubly-inductive construction:

- (i) $B_0 = Var(\phi)$ (B_0 is the set of atomic components of epistemic depth 0)
- (i') $D_0 = DNF(B_0)$ (D_0 is the set of disjunctive normal forms based on B_0)
- (ii) $B_k = \{L_a\psi : \psi \in D_{k-1}\}$
- (ii') $D_k = DNF(\bigcup_{l=0}^k B_l)$.

Notation B1. Let $\psi \in DNF(X)$ and $X \subseteq Y$. The expansion of ψ in $DNF(Y)$ is the formula obtained by the replacement of each conjunction $e_1\psi_1 \wedge \dots \wedge e_m\psi_m$ occurring in ψ by a disjunction of all possible conjunctions built from $e_1\psi_1 \wedge \dots \wedge e_m\psi_m$ by adding literals of atomic components in $Y - X$.

For instance, consider $X = \{p\}$ and $Y = \{p, q\}$: the DNF p is expanded in $(p \wedge q) \vee (p \wedge \neg q)$.

Fact B1.

- (i) $\forall k, \forall \psi \in D_k \cup B_k, dp(\psi) = k$.
- (ii) For each $\psi \in D_k$, in each $D_l, l > k$, there is a formula ψ' , which is equivalent to ψ

Proof. (i) is obvious; (ii) follows from the fact that any formula $\psi \in D_k$ can be expanded in ψ' from $D_l, l > k$ and that (by propositional reasoning) ψ and ψ' are equivalent. \square

It can be proven that $\mathfrak{B} = D_n$ is a finite base for $\mathcal{L}_{[\phi]}^L$. First, for each $\psi \in \mathcal{L}_{[\phi]}^L$, there is a formula ψ' in D_l s.t. $\vdash_{HM} \psi \leftrightarrow \psi'$ where $dp(\psi) = l$. Since ψ' can be expanded in a logically-equivalent formula $\psi'' \in D_n$, it is sufficient to conclude that for each $\psi \in \mathcal{L}_{[\phi]}^L$, there is an equivalent formula in the base.

- (i) $\psi := p$: ψ is obviously equivalent to some DNF in D_0 and $dp(p) = 0$.
- (ii) $\psi := (\chi_1 \wedge \chi_2)$: by the induction hypothesis, there is χ'_1 equivalent to χ_1 in $D_{dp(\chi_1)}$ and χ'_2 equivalent to χ_2 in $D_{dp(\chi_2)}$. Suppose w.l.o.g. that $dp(\chi_2) > dp(\chi_1)$ and, therefore, that $dp(\psi) = dp(\chi_2)$. Then, χ'_1 can be expanded in $\chi''_1 \in D_{dp(\chi_2)}$. Obviously, the disjunction of the conjunctions occurring both in χ''_1 and χ'_2 is in $D_{dp(\chi_2)}$ and equivalent to ψ .
- (iii) $\psi := L_a\chi$: by IH, there is χ' equivalent to χ in $D_{dp(\chi)}$. Note that $dp(\chi) < n = dp(\phi)$. By construction, $L_a(\chi') \in B_{dp(\chi)+1}$. Consequently, there will be in $D_{dp(\chi)+1}$ a DNF ψ' equivalent to $L_a(\chi')$. Since $dp(\chi) + 1 \leq n$, this DNF can be associated by expansion to a DNF in the base D_n . Furthermore, since $\vdash_{HM} \chi \leftrightarrow \chi'$ and $\vdash_{HM} L_a\chi' \leftrightarrow \psi'$, it follows by the rule of equivalence that $\vdash_{HM} L_a\chi \leftrightarrow \psi'$.

\square

There are changes needed to deal with unawareness. (1) First, the awareness operator A has to be included. This is not problematic given that for any formula $\psi, \vdash_{HM_U} A\psi \leftrightarrow \bigwedge_m A p_m$ where $Var(\psi) = \{p_1, \dots, p_m, \dots, p_M\}$. Consequently, the only modification is to include any formula $A p$ with $p \in Var(\phi)$

in B_0 . (2) With unawareness, it is no longer true that if $\vdash_{HM_U} \psi \leftrightarrow \chi$, then $\vdash_{HM_U} L_a \psi \leftrightarrow L_a \chi$. For instance, it is not true that $\vdash_{HM_U} L_a p \leftrightarrow L_a((p \wedge q) \vee (p \wedge \neg q))$: the agent may be unaware of q . Nevertheless, the rule holds restrictedly: under the assumption that $Var(\psi) = Var(\chi)$, then if $\vdash_{HM_U} \psi \leftrightarrow \chi$, then $\vdash_{HM_U} L_a \psi \leftrightarrow L_a \chi$ (RE_U). We can use this fact to make another change to the basis: instead of considering only the disjunctive normal forms built from the whole set $Var(\phi)$, we consider the disjunctive normal forms built from any non-empty subset $X \subseteq Var(\phi)$.

Definition B3. Let $X \subseteq Var(\phi)$;

- (i) $B_0^X = X \cup \{Ap : p \in X\}$
- (i') $D_0^X = DNF(B_0)$

- (ii) $B_k^X = \{L_a \psi : \psi \in D_{k-1}\}$ and
- (ii') $D_k^X = DNF(\bigcup_{l=0}^k B_l^X)$.

Fact B2.

- (i) $\forall k \leq n, \psi \in D_k^X$ where $X \subseteq Var(\phi)$, $dp(\psi) = k$.
- (ii) $\forall X \subseteq Var(\phi), \forall \psi \in D_k^X, \forall D_l^X, l > k$, there is a formula $\psi' \in D_l^X$, which is equivalent to ψ .
- (iii) $\forall X \subseteq Y \subseteq Var(\phi)$, if $\psi \in D_k^X$, then there is a formula ψ' , which is equivalent to ψ in D_k^Y .
- (iv) $\forall X \subseteq Var(\phi), \forall \psi \in D_k^X, Var(\psi) = X$.

Proof. (i)–(ii) are similar to the classical case; (iii) is straightforwardly implied by Clause (ii') of Definition 12; (iv) is obvious. \square

We are now ready to prove that $\mathfrak{B} = \bigcup_{X \subseteq Var(\phi)} D_n^X$ is a basis for $\mathcal{L}_{[\phi]}^{LA}$. We will actually show that for any $\psi \in \mathcal{L}_{[\phi]}^{LA}$ with $dp(\psi) = k$, there are $X \subseteq Var(\phi)$ and $\psi' \in D_k^X$ s.t. $\vdash_{HM_U} \psi \leftrightarrow \psi'$ and $dp(\psi) = dp(\psi') = k$ and $Var(\psi) = Var(\psi')$ (Induction Hypothesis, IH).

- (i) $\psi := p$: ψ is obviously equivalent to some DNF ψ' in $D_0^{\{p\}}$. Clearly, $dp(\psi) = dp(\psi')$ and $Var(\psi) = Var(\psi')$.
- (ii) $\psi := (\chi_1 \wedge \chi_2)$: by IH,
 - there is χ'_1 s.t. $\vdash_{HM_U} \chi_1 \leftrightarrow \chi'_1$ and $Var(\chi_1) = Var(\chi'_1) = X_1$ and $\chi'_1 \in D_{dp(\chi_1)}^{X_1}$
 - there is χ'_2 s.t. $\vdash_{HM_U} \chi_2 \leftrightarrow \chi'_2$ and $Var(\chi_2) = Var(\chi'_2) = X_2$ and $\chi'_2 \in D_{dp(\chi_2)}^{X_2}$
 Let us consider $X' = X_1 \cup X_2$ and suppose without loss of generality that $dp(\chi_2) > dp(\chi_1)$. One may expand χ'_1 from $D_{dp(\chi_1)}^{X_1}$ to $D_{dp(\chi_1)}^{X'}$ and expand the resulting DNF to $\chi''_1 \in D_{dp(\chi_2)}^{X'}$. On the other hand, χ'_2 may be expanded to $\chi''_2 \in D_{dp(\chi_2)}^{X'}$. ψ' is the disjunction of the conjunctions common to χ''_1 and χ''_2 . Obviously, $dp(\psi) = dp(\psi')$ and $Var(\psi) = Var(\psi')$.
- (iii) $\psi := A\chi$: by IH, there is χ' equivalent to χ in $D_{dp(\chi)}^X$ with $Var(\chi) = Var(\chi')$. $A\chi'$ is equivalent to $\bigwedge_m Ap_m$ where $Var(\chi') = \{p_1, \dots, p_m, \dots, p_M\}$. Each Ap_m is in B_0^X , so by expansion in $D_{dp(\chi)}^X$, there is a DNF equivalent to it and, therefore, a DNF equivalent to $\bigwedge_m Ap_m$.
- (iv) $\psi := L_a \chi$: by IH, there is χ' equivalent to χ in $D_{dp(\chi)}^X$ with $dp(\chi) = dp(\chi')$ and $Var(\chi) = Var(\chi')$. Note that $dp(\chi) < n = dp(\phi)$. By construction, $L_a(\chi') \in B_{dp(\chi)+1}^X$. Consequently, there will be in $D_{dp(\chi)+1}^X$ a DNF ψ' logically equivalent to $L_a(\chi')$. Since $dp(\chi) + 1 \leq n$, there will be in the base a formula ψ'' logically equivalent to ψ' . Furthermore, since $\vdash_{HM_U} \chi \leftrightarrow \chi'$ and $Var(\chi) = Var(\chi')$ and $\vdash_{HM_U} L_a \chi' \leftrightarrow \psi''$, it follows that $\vdash_{HM_U} L_a \chi \leftrightarrow \psi''$.

We will now build: (1) the objective state space; (2) the subjective states spaces and the projection ρ ; and (3) the probability distributions.

- (1) The objective state space:

The objective states of the ϕ -canonical structure are the intersections of the maximally-consistent sets of formulas of the language $\mathcal{L}^{LA}(At)$ and the restricted language $\mathcal{L}_{[\phi]}^{LA}$:

$$S^\phi = \{\Gamma \cap \mathcal{L}_{[\phi]}^{LA} : \Gamma \text{ is a maximal } HM_U\text{-consistent set}\}$$

First, let us notice that the system HM_U is a “modal logic” ([38], p. 191): a set of formulas (1) that contains every propositional tautologies; (2) such that the Lindenbaum lemma holds.

Definition B4.

- (i) A formula ϕ is deducible from a set of formulas Γ , symbolized $\Gamma \vdash_{HM_U} \phi$, if there exists some formulas ψ_1, \dots, ψ_n in Γ s.t. $\vdash_{HM_U} (\psi_1 \wedge \dots \wedge \psi_n) \rightarrow \phi$.
- (ii) A set of formulas Γ is HM_U -consistent if it is false that $\Gamma \vdash_{HM_U} \perp$.
- (iii) A set of formulas Γ is maximally HM_U -consistent if (1) it is HM_U -consistent and (2) if it is not included in a HM_U -consistent set of formulas.

Lemma B1 (Lindenbaum Lemma). *If Γ is a set of HM_U -consistent formulas, then there exists an extension Γ^+ of Γ that is maximally HM_U -consistent.*

Proof. See, for instance, [38] (p.199). \square

Notation B2. For each formula $\psi \in \mathcal{L}_{[\phi]}^{LA}$, let us note $[\psi] = \{s \in S^\phi : \psi \in s\}$

Lemma B2. *The set S^ϕ is finite.*

Proof. This Lemma is a consequence of the fact that $\mathcal{L}_{[\phi]}^{LA}$ is finitely generated.

- (a) Let us say that two sets of formulas are Δ -equivalent if they agree on each formula that belongs to Δ . S^ϕ identifies the maximal HM_U -consistent sets of formulas that are $\mathcal{L}_{[\phi]}^{LA}$ -equivalent. S^ϕ is infinite iff there are infinitely many maximal HM_U -consistent sets of formulas that are not pairwise $\mathcal{L}_{[\phi]}^{LA}$ -equivalent.
- (b) If \mathfrak{B} is a base for $\mathcal{L}_{[\phi]}^{LA}$, then two sets of formulas are $\mathcal{L}_{[\phi]}^{LA}$ -equivalent iff they are \mathfrak{B} -equivalent. Suppose that Δ_1 and Δ_2 are not $\mathcal{L}_{[\phi]}^{LA}$ -equivalent. This means w.l.o.g. that there is a formula ψ s.t. (i) $\psi \in \Delta_1$, (ii) $\psi \notin \Delta_2$ and (iii) $\psi \in \mathcal{L}_{[\phi]}^{LA}$. Let $\psi' \in \mathfrak{B}$ be a formula s.t. $\vdash_{HM_U} \psi \leftrightarrow \psi'$. Clearly, $\psi' \in \Delta_1$ and $\psi' \in \mathcal{L}_{[\phi]}^{LA}$ and $\neg\psi' \in \Delta_2$. Therefore, Δ_1 and Δ_2 are not \mathfrak{B} -equivalent. The other direction is obvious.
- (c) Since \mathfrak{B} is finite, there are only finitely many maximal HM_U -consistent sets of formulas that are not pairwise \mathfrak{B} -equivalent. Therefore, S^ϕ is finite.

\square

(2) The subjective state spaces and the projection $\rho(\cdot)$:

As it might be expected, the subjective state associated with an objective state $\Gamma \in S^\phi$ will be determined by the formulas that the agent is aware of in Γ .

Definition B5. *For any set of formulas Γ , let $Var(\Gamma)$ be the set of atomic formulas that occur in the formulas that belong to Γ . For any $\Gamma \in S^\phi$, let:*

- (i) $A^+(\Gamma) = \{\psi : A\psi \in \Gamma\}$ and $A^-(\Gamma) = \{\psi : \neg A\psi \in \Gamma\}$
- (ii) $a^+(\Gamma) = \{p \in Var(\mathcal{L}_{[\phi]}^{LA}) : Ap \in \Gamma\}$ and $a^-(\Gamma) = \{p \in Var(\mathcal{L}_{[\phi]}^{LA}) : \neg Ap \in \Gamma\}$.

Lemma B3. *Let $\Gamma \in S^\phi$.*

- (i) $A^+(\Gamma) = \mathcal{L}_{[\phi]}^{LA}(a^+(\Gamma))$

$$(ii) A^+(\Gamma) \cup A^-(\Gamma) = \mathcal{L}_{[\phi]}^{LA}$$

Proof. (i) Follows from (A1)-(A4_L); (ii) follows from (i) and the fact that since Γ comes from a maximal consistent set, $\neg\psi \in \Gamma$ iff $\psi \notin \Gamma$. \square

One may group the sets that have the same awareness profile into equivalence classes: $|\Gamma|_a = \{\Delta \in S^\phi : a^+(\Delta) = a^+(\Gamma)\}$. The sets that belong to the same equivalence class $|\Gamma|_a$ will be mapped in the same subjective state space $S'_{|\Gamma|_a}$. We are now ready to define the projection ρ and these subjective states.

Definition B6. The projection $\rho : S^\phi \rightarrow \bigcup_{\Gamma \in S^\phi} S'_{|\Gamma|_a}$ is defined by:

$$\rho(\Gamma) = \Gamma \cap A^+(\Gamma)$$

where $S'_{|\Gamma|_a} = \{\Delta \cap A^+(\Gamma) \cap \mathcal{L}_{[\phi]}^{LA} : \Delta \text{ is a maximal } HM_U\text{-consistent set and } a^+(\Delta) = a^+(\Gamma)\}$.

Note that in the particular case where the agent is unaware of every formula, $A^+(\Gamma) = \emptyset$. Therefore, each objective state where the agent is unaware of every formula will be projected in the same subjective state $\emptyset \in S'_\emptyset = \{\emptyset\}$. More importantly, one can check that ρ is an onto map: suppose that $\Lambda \in S'_{|\Gamma|_a}$ where $\Gamma \in S^\phi$. By definition, for some Δ (a maximal HM_U -consistent set), $\Lambda = \Delta \cap A^+(\Gamma) \cap \mathcal{L}_{[\phi]}^{LA}$ and $a^+(\Delta) = a^+(\Gamma)$. As a consequence, $A^+(\Delta) = A^+(\Gamma)$, and therefore, $\Lambda = \Delta \cap A^+(\Delta) \cap \mathcal{L}_{[\phi]}^{LA}$. Hence, $\Lambda = \rho(\Delta \cap \mathcal{L}_{[\phi]}^{LA})$. One can show also the following lemma.

Lemma B4.

- (i) For each $\Gamma \in S^\phi$, $S'_{|\Gamma|_a}$ is finite.
- (ii) For each subset $E \subseteq S'_{|\Gamma|_a}$, there is $\psi \in A^+(\Gamma) \cap \mathcal{L}_{[\phi]}^{LA}$ s.t. $E = [\psi]_{S'_{|\Gamma|_a}}$ where $[\psi]_{S'_{|\Gamma|_a}}$ denotes the set of states of $S'_{|\Gamma|_a}$ to which ψ belongs.
- (iii) For all $\psi_1, \psi_2 \in A^+(\Gamma) \cap \mathcal{L}_{[\phi]}^{LA}$, $[\psi_1]_{S'_{|\Gamma|_a}} \subseteq [\psi_2]_{S'_{|\Gamma|_a}}$ iff $\vdash_{HM_U} \psi_1 \rightarrow \psi_2$

Proof. (i) Follows trivially since the objective state space is already finite; (ii) let us pick a finite base \mathfrak{B}_Γ for $A^+(\Gamma) \cap \mathcal{L}_{[\phi]}^{LA}$. For each element β of this base and each $\Delta \in S'_{|\Gamma|_a}$, either $\beta \in \Delta$ or $\neg\beta \in \Delta$. Two sets Δ and $\Delta' \in S'_{|\Gamma|_a}$ differ at least by one such formula of \mathfrak{B}_Γ . Let $C(\Delta) = \bigwedge_m e_m \beta_m$ where $\beta_m \in \mathfrak{B}_\Gamma$ and e_m is a blank if $\beta_m \in \Delta$ and \neg if $\beta_m \notin \Delta$. For two distinct sets Δ and Δ' , $C(\Delta) \neq C(\Delta')$. For each event $E \subseteq S'_{|\Gamma|_a}$, one can therefore consider the disjunction $\bigvee_k C(\Delta_k)$ for each $\Delta_k \in E$. Such a formula belongs to each Δ_k and only to these Δ_k . (iii) (\Rightarrow). For each formula $\psi \in A^+(\Gamma) \cap \mathcal{L}_{[\phi]}^{LA}$ and each $\Delta \in S'_{|\Gamma|_a}$, $\neg\psi \in \Delta$ iff $\psi \notin \Delta$. Therefore, there are two possibilities for any Δ : either $\psi \in \Delta$ or $\neg\psi \in \Delta$. (a) If $\psi_1 \in \Delta$, then by hypothesis $\psi_2 \in \Delta$ and given the construction of the language, $\neg\psi_1 \vee \psi_2 \in \Delta$, hence $\psi_1 \rightarrow \psi_2 \in \Delta$. (b) If $\psi_1 \notin \Delta$, then $\neg\psi_1 \in \Delta$, hence $\psi_1 \rightarrow \psi_2 \in \Delta$. This implies that for any Δ , $\psi_1 \rightarrow \psi_2 \in \Delta$. Given the definition of $S'_{|\Gamma|_a}$ and the properties of maximal consistent sets, this implies that $\vdash_{HM_U} \psi_1 \rightarrow \psi_2$. (\Leftarrow). Given the construction of the language, if $\psi_1, \psi_2 \in A^+(\Gamma) \cap \mathcal{L}_{[\phi]}^{LA}$, then $\psi_1 \rightarrow \psi_2 \in A^+(\Gamma) \cap \mathcal{L}_{[\phi]}^{LA}$. Since $\vdash_{HM_U} \psi_1 \rightarrow \psi_2$, for each Δ , $\psi_1 \rightarrow \psi_2 \in \Delta$. If $\psi_1 \in \Delta$, clearly $\psi_2 \in \Delta$, as well. Therefore, $[\psi_1]_{S'_{|\Gamma|_a}} \subseteq [\psi_2]_{S'_{|\Gamma|_a}}$. \square

(3) The probability distributions:

Definition B7. For $\Gamma \in S^\phi$ and $\psi \in \mathcal{L}_{[\phi]}^{LA}$, let:

- $\tilde{a} = \max\{a : L_a \psi \in \Gamma\}$
- $\tilde{b} = \min\{b : M_b \psi \in \Gamma\}$

In the classical case [19], \tilde{a} and \tilde{b} are always defined. This is not so in our structure with unawareness: if the agent is not aware of ψ , no formula $L_a\psi$ will be true because of $(A0_U) A\psi \leftrightarrow L_0\psi$. Given (A1) and (DefM), one can derive:

$$\vdash_{HM_U} A\psi \leftrightarrow M_1\psi$$

The construction of the language implies that for any Γ , $A\psi \in \Gamma$ iff $L_0\psi \in \Gamma$ iff $M_1\psi \in \Gamma$. Therefore, \tilde{a} and \tilde{b} are defined iff $A\psi \in \Gamma$.

Lemma B5. *Let us suppose that $A\psi \in \Gamma$.*

- (i) $\forall c \in Q(\phi)$, $c \leq \tilde{a}$ implies $L_c\psi \in \Gamma$, and $c \geq \tilde{b}$ implies $M_c\psi \in \Gamma$
- (ii) There are only two cases: (i) either $\tilde{a} = \tilde{b}$ and $E_{\tilde{a}}\psi \in \Gamma$ while $E_c\psi \notin \Gamma$ for $c \neq \tilde{a}$, (ii) or $\tilde{a} < \tilde{b}$ and $E_c\psi \notin \Gamma$ for any $c \in Q(\phi)$.
- (iii) $\tilde{b} - \tilde{a} \leq \frac{1}{q}$ (where q is the common denominator to the indexes)

Proof. See [19]; the modifications are obvious. \square

Definition B8. *Given $\Gamma \in S^\phi$ and $\psi \in \mathcal{L}_{[\phi]}^{LA}$, if $A\psi \in \Gamma$, let*

I_ψ^Γ be either $\{\tilde{a}\}$ if $\tilde{a} = \tilde{b}$ or (\tilde{a}, \tilde{b}) if $\tilde{a} < \tilde{b}$.

Lemma A.5 in [19] can be adapted to show that for each $S'_{|\Gamma|_a}$ and $\Gamma \in S'_{|\Gamma|_a}$, there is a probability distribution $P'_{|\Gamma|_a}(\Gamma)$ on $2^{S'_{|\Gamma|_a}}$, such that

$$(C) \text{ for all } \psi \in \mathcal{L}_{[\phi]}^{LA} \text{ if } A\psi \in \Gamma, P'_{|\Gamma|_a}(\Gamma)([\psi]_{S'_{|\Gamma|_a}}) \in I_\psi^\Gamma.$$

The proof in [19] relies on a theorem by Rockafellar that can be used because of the inference rule (B). It would be tedious to adapt the proof here. One comment is nonetheless important. In our axiom system HM_U , the inference rule holds under a restricted form (B_U) . Therefore, one could wonder whether this will not preclude adapting the original proof, which relies on the unrestricted version (B). It turns out that the answer is negative. The reason is that the formulas involved in the application of (B) are only representatives for each subset of the state space. We have previously shown how to build these formulas in our case, and they are such that the agent is necessarily aware of them. Therefore, the restriction present in (B_U) does not play any role, and we may define the ϕ -canonical structure as follows.

Definition B9. *The ϕ -canonical structure is the GSPS $\mathcal{M}_\phi = (S^\phi, S^{\phi'}, (2^{S'_{|\Gamma|_a}})_{\Gamma \in S^\phi}, \pi_\phi, (P'_{|\Gamma|_a})_{\Gamma \in S^\phi})$ where:*

- (i) $S^\phi = \{\Gamma \cap \mathcal{L}_{[\phi]}^{LA} : \Gamma \text{ is a maximal } HM_U\text{-consistent set}\}$
- (ii) $S^{\phi'} = \bigcup_{\Gamma \in S^\phi} S'_{|\Gamma|_a}$ where $S'_{|\Gamma|_a} = \{\Delta \cap A^+(\Gamma) \cap \mathcal{L}_{[\phi]}^{LA} : \Delta \text{ is a maximal } HM_U\text{-consistent set, and } a(\Delta) = a(\Gamma)\}$
- (iii) for each $\Gamma \in S^\phi$, $\rho(\Gamma) = \Gamma \cap A^+(\Gamma)$
- (iv) for all state $\Gamma \in S^\phi \cup S^{\phi'}$ and atomic formula $p \in At$, $\pi_\phi(p, \Gamma) = 1$ iff $p \in \Gamma$
- (v) for $\Gamma \in S^\phi$, $P'_{|\Gamma|_a}$ is a probability distribution on $2^{S'_{|\Gamma|_a}}$ satisfying Condition (C)²⁴.

We are now ready to state the crucial truth lemma.

Lemma B6 (Truth lemma). *For every $\Gamma \in S^\phi$ and every $\psi \in \mathcal{L}_{[\phi]}^{LA}$,*

²⁴ In the particular case where $A^+(\Gamma) = \emptyset$, the probability assigns maximal weight to the only state of S'_\emptyset .

$$\mathcal{M}_\phi, \Gamma \models \psi \text{ iff } \psi \in \Gamma$$

Proof. The proof proceeds by induction on the complexity of the formula.

- $\psi := p$; following directly from the definition of π_ϕ .
- $\psi := \neg\chi$. Since Γ is a standard state $\mathcal{M}_\phi, \Gamma \models \neg\chi$ iff $\mathcal{M}_\phi, \Gamma \not\models \chi$ iff (by IH) $\chi \notin \Gamma$. We shall show that $\chi \notin \Gamma$ iff $\neg\chi \in \Gamma$. (\Rightarrow) Let us suppose that $\chi \notin \Gamma$; χ is in $\mathcal{L}_{[\phi]}^{LA}$; hence, given the properties of maximally-consistent sets, $\neg\chi \in \Gamma^+$ where Γ^+ is the extension of Γ to $\mathcal{L}^{LA}(At)$ (the whole language). Additionally, since $\Gamma = \Gamma^+ \cap \mathcal{L}_{[\phi]}^{LA}$, $\neg\chi \in \Gamma$. (\Leftarrow) Let us suppose that $\neg\chi \in \Gamma$. Γ is coherent, therefore $\chi \notin \Gamma$.
- $\psi := \psi_1 \wedge \psi_2$. (\Rightarrow). Let us assume that $\mathcal{M}_\phi, \Gamma \models \psi_1 \wedge \psi_2$. Then, $\mathcal{M}_\phi, \Gamma \models \psi_1$ and $\mathcal{M}_\phi, \Gamma \models \psi_2$. By IH, this implies that $\psi_1 \in \Gamma$ and $\psi_2 \in \Gamma$. Given the properties of maximally-consistent sets, this implies in turn that $\psi_1 \wedge \psi_2 \in \Gamma$. (\Leftarrow). Let us assume that $\psi_1 \wedge \psi_2 \in \Gamma$. Given the properties of maximally-consistent sets, this implies that $\psi_1 \in \Gamma$ and $\psi_2 \in \Gamma$ and, therefore, by IH, that $\mathcal{M}_\phi, \Gamma \models \psi_1$ and $\mathcal{M}_\phi, \Gamma \models \psi_2$.
- $\psi := A\chi$. We know that in any GSPS \mathcal{M} , if $s' = \rho(s) \in S'_X$ for some $s \in S$, then $\mathcal{M}, s' \Downarrow \chi$ iff $\chi \in \mathcal{L}^{LA}(X)$. In our case, $s = \Gamma$, $s' = \rho(\Gamma)$ and $X = a^+(\Gamma)$. Therefore, $\mathcal{M}_\phi, \Gamma \models A\chi$ iff $\chi \in \mathcal{L}^{LA}(a^+(\Gamma))$. However, given that $A\chi \in \mathcal{L}_{[\phi]}^{LA}$, $\chi \in \mathcal{L}^{LA}(a^+(\Gamma))$ iff $A\chi \in \Gamma$.
- $\psi := L_a\chi$. By definition $\mathcal{M}_\phi, \Gamma \models L_a\chi$ iff $P_{|\rho(\Gamma)|_a}(\Gamma)([[\chi]]) \geq a$ and $\mathcal{M}_\phi, \rho(\Gamma) \Downarrow \chi$. (\Leftarrow) Let us suppose that $P_{|\rho(\Gamma)|_a}(\Gamma)([[\chi]]) \geq a$ and $\mathcal{M}_\phi, \rho(\Gamma) \Downarrow \chi$. Hence, \tilde{a} is well defined. It is clear that $\tilde{a} \geq a$ given our definition of $P_{|\rho(\Gamma)|_a}(\Gamma)$. It is easy to see that $\vdash_{HM_U} L_a\psi \rightarrow L_b\psi$ for $b \leq a$. As a consequence, $L_a\psi \in \Gamma$. (\Rightarrow) Let us suppose that $L_a\psi \in \Gamma$. This implies that $A\psi \in \Gamma$ and, therefore, that $\mathcal{M}_\phi, \rho(\Gamma) \Downarrow \chi$. By construction, $a \leq \tilde{a}$, and therefore, $P_{|\rho(\Gamma)|_a}(\Gamma)([[\chi]]) \geq a$. Hence, $\mathcal{M}_\phi, \Gamma \models L_a\chi$.

□

Proof.

- If $\phi := p$, then $\mathcal{M}, s' \models p$ or $\mathcal{M}, s' \models \neg p$
iff $(p \in \mathcal{L}^{LA}(X) \text{ and } \pi^*(s', p) = 1) \text{ or } (p \in \mathcal{L}^{LA}(X) \text{ and } \pi^*(s', p) = 0)$
iff $p \in \mathcal{L}^{BA}(X)$
- If $\phi := \neg\psi$, then $\mathcal{M}, s' \models \phi$ or $\mathcal{M}, s' \models \neg\phi$
iff $\mathcal{M}, s' \models \neg\psi$ or $\mathcal{M}, s' \models \neg\neg\psi$
iff $(\psi \in \mathcal{L}^{LA}(X) \text{ and } \mathcal{M}, s' \not\models \psi) \text{ or } (\neg\psi \in \mathcal{L}^{LA}(X) \text{ and } \mathcal{M}, s' \not\models \neg\psi)$
iff $(\psi \in \mathcal{L}^{LA}(X) \text{ and } \mathcal{M}, s' \not\models \psi) \text{ or } (\psi \in \mathcal{L}^{LA}(X) \text{ and } \mathcal{M}, s' \models \psi)$
iff $\psi \in \mathcal{L}^{LA}(X)$
iff $\neg\psi \in \mathcal{L}^{LA}(X)$
- If $\phi := \psi_1 \wedge \psi_2$, then $\mathcal{M}, s' \models \phi$ or $\mathcal{M}, s' \models \neg\phi$
iff $(\mathcal{M}, s' \models \psi_1 \text{ and } \mathcal{M}, s' \models \psi_2) \text{ or } (\psi_1 \wedge \psi_2 \in \mathcal{L}^{LA}(X) \text{ and } (\mathcal{M}, s' \not\models \psi_1 \text{ or } \mathcal{M}, s' \not\models \psi_2))$
iff by the induction hypothesis $(\psi_1 \wedge \psi_2 \in \mathcal{L}^{LA}(X) \text{ and } \mathcal{M}, s' \models \psi_1 \text{ and } \mathcal{M}, s' \models \psi_2) \text{ or } (\psi_1 \wedge \psi_2 \in \mathcal{L}^{LA}(X) \text{ and not } (\mathcal{M}, s' \models \psi_1 \text{ and } \mathcal{M}, s' \models \psi_2))$
iff $\psi_1 \wedge \psi_2 \in \mathcal{L}^{LA}(X)$

- if $\phi := B\psi$, then $\mathcal{M}, s' \models \phi$ or $\mathcal{M}, s' \models \neg\phi$
 iff (for each $t^* \in R^*(s')$, $\mathcal{M}, t^* \models \phi$) or ($B\psi \in \mathcal{L}^{LA}(X)$) and $\mathcal{M}, s' \not\models B\psi$
 iff, by the induction hypothesis and since each $t^* \in R^*(s')$ belongs to $S_X - (B\psi \in \mathcal{L}^{LA}(X))$
 $\mathcal{M}, s' \models B\psi$ or ($B\psi \in \mathcal{L}^{LA}(X)$ and $\mathcal{M}, s' \not\models B\psi$)
 iff $B\psi \in \mathcal{L}^{LA}(X)$
 - If $\phi := A\psi$, then $\mathcal{M}, s' \models \phi$ or $\mathcal{M}, s' \models \neg\phi$
 iff ($\mathcal{M}, s' \Downarrow \psi$) or ($A\psi \in \mathcal{L}^{LA}(X)$ and not $\mathcal{M}, s' \Downarrow \psi$, impossible given the induction hypothesis)
 iff ($\mathcal{M}, s' \Downarrow \psi$)
 iff $\psi \in \mathcal{L}^{LA}(X)$ (by the induction hypothesis)
 iff $A\psi \in \mathcal{L}^{LA}(X)$
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